ASM Programming: Ch. 2-3

Outline

- Structured programming process
- Program structures
- Looping concept
- Looping constructs
- ASM code examples
- Branch instructions
- Relative address mode
- BRSET/BRCLR instructions
- Instruction timing
- Delay loops
- ASM loop/branch examples



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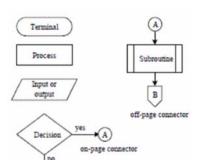
Structured Programming Process

1. Define Problem

- -what exactly should your program do?
- 2. Plan Solution = Develop Algorithm
 - -sequence of computational or logical step to transform given inputs to desired outputs *Options*
- top-down: start with goals, work down to required outputs
- <u>bottom-up</u>: start with required *outputs*, work up to *goals*
- <u>divide & conquer</u>
 - use top-down to break into tasks
 - use bottom-up to solve each task
 - best method for large programs

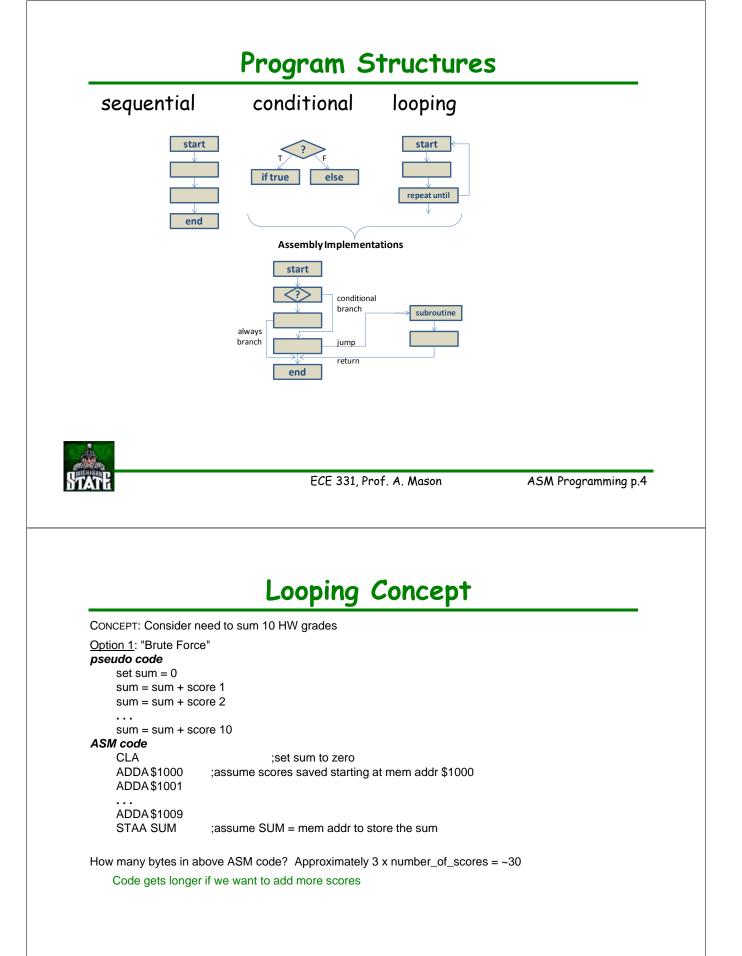
Approaches for algorithm planning

- <u>Pseudocode</u>: non-code list of tasks (example will follow)
- <u>Flowcharts</u>: graphical map of algorithm flow (example follow)
- 3. Code Algorithm
 - -ASM, C, higher level language
- 4. Simulate & Debug
- 5. Test in Hardware & Debug







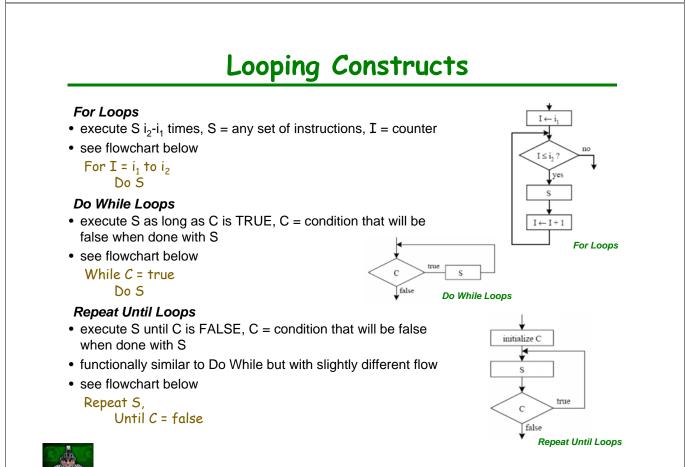




Looping Concept

<u>Option 2</u> : Use a Loop pseudo code set sum = 0 set count = 0 while count < 10 sum = sum + score(count) increment count (loop back to "while") ASM code See "Simple Looping Code" Can you identify the loop in this program?	.ASM co ; Loop ex ; Sums 10 values to	0 values fr sum are s	<i>text</i> Ch. 3 notes	
How many bytes of code in the loop? 21, for an infinite number of scores	DONE SUM	STAB EQU END	SUM \$4400	;store result
LST compiled file, text Line Addr Op Code Label Mnemonic Operand Not 1: ; Loop example for Ch. 3 note 2: ; Sums 10 values from memory 3: ; assumes 10 values to sum ar 4: ; assumes prior sum stored at	s by A.Mason. Mar & store result in e stored at \$1000	SUM		

1:		; Loop	example :	for Ch.	3 notes by A.Mason. Mar 09	
2:		; Sums	10 values	s from m	memory & store result in SUM	
3:		; assum	nes 10 va	lues to	sum are stored at \$1000	
4:		; assum	mes prior	sum sto	ored at SUM	
5:	=00004000		ORG	\$4000		
6:	4000 CE 1000		LDX	#\$1000	;set x to fist memory addr	
7:	4003 F6 4400		LDAB	SUM	;load staring sum into accB	
8:	4006 86 00		LDAA	#\$00	; initialize counter to 0	
9:	4008 81 0A	CHECK	CMPA	#\$0A	; ?added all 10?	.S19 machine code record, binary
10:	400A 27 06		BEQ	DONE	; if yes, done	S0030000FC
11:	400C EB 00		ADDB	0,X	; if no, add # to SUM	S1134000CE1000F644008600810A2706EB00084221
12:	400E 08		INX		; increment IX	S108401020F67B4400D2
13:	400F 42		INCA		; increment counter	S9030000FC
14:	4010 20 F6		BRA	CHECK	;repeat loop	
15:	4012 7B 4400	DONE	STAB	SUM	;store result	
16:	=00004400	SUM	EQU	\$4400		
17:			END			ASM Programming p.6





ASM Examples

Discuss ASM Examples in HO_5, pg. 4-6

Simple Arithmetic ASM Program Examples

- Example 1: Write a program to add the numbers stored at memory locations \$800, \$801, and \$802, and store the sum at memory location \$900.
- Example 2: Write a program to subtract the contents of the memory location at \$805 from the sum of the memory locations at \$800 and \$802, and store the result at the memory location \$900.
- Example 3: Write a program to add two 16-bit numbers that are stored at \$800~\$801 and \$802~\$803, and store the sum at \$900~\$901.
- Example 4: Write a program to subtract 5 from four memory locations at \$800, \$801, \$802, and \$803.
- Assembly and Execution Example (Chapter 2)



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More ASM Examples

• See "Example ASM Code.txt

Simple AND example (A.Mason Feb 10) Function: AND value in accA with value in program memory and store to a Address mode for each instruction shown in comments version1: no labels URG \$4000 LDAA #4 ;immediate

\$4000 #\$0F #\$F0 \$5550 ; immediate STAA ;extended SWI

FSimple AND example (A.Mason Feb 10) FFunction: AND value in accA with value in memory and store to mem Address mode for each instruction shown in comments pression2: uses labels and equates, different values than version1 OCG \$4000 Mont
 Labels a

 ORG
 \$4000

 EQU
 \$AA

 EQU
 \$55

 EQU
 \$550

 LDAA
 #55
N11m2 _____ ____1+ \$5550 #Numl #Num2 Result ;immediate ;immediate ANDA STAA ;extended END

le (A.Mason Feb 10) Soundier And Example (Alason Feb 10) Founction: NAN value in accA with value in memory and store to mu Address mode for each instruction shown in comments version: uses data storage and gets operands from memory ORG \$4000 ORG LDAA ANDA STAA SWI \$4000 Numl Num2 Result ;stop ;extended ;extended ;extended ige ORG \$5550 \$FA \$5F \$00 ;initialize result byte [5552] to 00 ORG FCB FCB FCB Num2

;ECE331	1 Simple	e Summation	(A.Mason Feb 10)
;add se	eries o	f numbers be	ginning at 'nums'
			et by value in 'count'
			avoid overflow; 'count' must be < 10
	res inde	exed address	ing; to simplify code, numbers added in reverse
order			
		org	\$4000
		ldab	count ;# of numbers to add
	decb		;# of additions is 1 less than # of numbers
	ldx	#nums	;addr for nums
	ldaa	sum	;load starting sum
again	adda	b,x	
	decb		;decrement counter
	bpl	again	;if b >= zero, add another
	staa	sum	
		swi	
		org	\$6000
count		fcb \$0	13
nums		fcb	1,2,3,4,5,6,7,8,9,10
sum		fcb	<pre>\$F0 ;previous sum -random value</pre>
		end	

; ECE331 Example of SET/CLR Bit program will read data and set lower nibble to %0011 until odd value is read ORG \$4000 #00 ; initialize index offset LDAA byte LDX #DATA ; initialize index reference register ;end if odd value BRSET A,X,\$01,ODD TOP BSET A,X,%00000011 A,X,%00001100 BCLR

A,X

TOP

\$AA

\$6000 \$EE, \$DC, \$D0, \$F4 \$80, \$00, \$55, \$22

LDAB

INCA

BRA

SWI

OPC

FCB FCB FCB



ODD

DATA

data storage

:end

Branch Instructions

- Function of branch instructions
 - -implement loops; IF-THEN-ELSE constructs
 - -alter execution order of linear instructions
- 68HC12 has 2 types of branches
 - -conditional
 - -unconditional
 - -both use Relative Addr. Mode

Unconditional Branches

- -does not depend on any conditions; will always occur
- BRA ;branch always
 - 8-bit signed offset
 - distance to next instr. +127 to -128
- -EXAMPLE:

– best to use labels & let assembler determine relative offset



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Branch Instructions II

- More unconditional branches
 - -LBRA ;long branch always, 16-bit signed offset
 - JMP ;jump, 16-bit extended/indexed addr.
 - not relative like BRA; specific 16-bit addr.
 - less compact & more clocks than BRA

Conditional Branch

- -branch only if a condition is true
 - conditions derived from CCR flags
 - previous instr. must set proper CCR flags
 - can use a Test/Compare instruction to set flags without changing memory values
- Bxx ; conditional branches with 8-bit signed offset
- LBxx ;conditional long branches with 16-bit signed offset
- -all conditional branches use Relative Addr. Mode
- best to use labels & let assembler determine relative offset operand

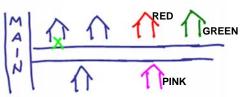


AMS Branch Instructions

	Uncondit	ional Branches		
Mnemonic	Function	Comment		
BRN	branch never	useful as delay		
BRA	branch always	8-bit signed offs	set	Example:
	Simple Con	ditional Branches		If all flags = 0 after
Mnemonic	"if" Function	Condition = Flag	g	an instruction, which
BCC	carry clear	C=0	-	branches would be
BCS	carry set	C=1		taken?
BEQ	equal	Z=1		
BNE	not equal	Z=0		
BMI	minus	N=I		
BPI	plus	N=0		
BVS	overflow set	V=1		
BVC	overflow clear	V=0		
	Unsigned Co	nditional Branches	3	
Mnemonic	"if" Function	Condition	Flags	
BHS	higher or same	$r \ge 0$	C=0	
BHI	higher than	r > 0	C+Z=0	
BLS	lower or same	r ≤ 0	C+Z=1	
BLO	lower than	r < 0	C=1	
	Signed Con	ditional Branches		
Mnemonic	"if" Function	Condition	Flags	
BGE	greater or equal	$r \ge 0$	N⊕V=0	
BGT	greater than	r > 0	Z+N⊕V=0	
BLE	less or equal	r ≤ 0	Z+N⊕V=1	
BLT	less than	r < 0	N⊕V=1	
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"Relative" Concept

• How can we define the address of the red house?



Absolute Address

- -3rd House on the left from Main Street
- Relative Address
 - -locate by offset from reference
 - -<u>references</u> offsets offset + reference
 - Green House
 - Pink House
 - 'X'



Relative Address Mode

• Function

- -adjust PC by value in operand
- -operand can be set by program or by assembler
 - assembler can calculate "distance" to a Label on line to branch to
- -only used in Branch instr.

• EXAMPLES



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• Offset	in Relative Add	dr. Mode sp	ecifies how many instruction
		•	PC must be moved to get
	urrent PC value	,	6
-PC(ne	ext) = PC(current)	+ Offset ;off	set is signed (+/-)
• Numbe	er of Bytes per	Instruction	
	of instr. op-code a		/tes
-easy t	to see in .LST file		
• EXAM	PLE: What is t	he hex valu	e of ??
Program		op-codes	#bytes
-	LDDA #10	BCOA	2
STUK	INCA	42	1
	STAA \$10B5	A7 01 B5	3
	TSTA	97	1
	BNE STUK	26 ??	2

BRCLR / BRSET



- -combine test and branch instr.
- -can use direct, extended, or indexed addr. mode
- Format
 - BRCLR/SET mem mask label
 - mem = addr. of data to test
 - mask = bits to check (set to 1 if checked)
 - label = branch location
 - -BRCLR: are 1 bits in mask = 0 @ mem?
 - if yes, branch to label
 - -BRSET: are 1 bits in mask = 1 @ mem?
 - if yes, branch to label

• EXAMPLE BRCLR 0,X %11110000, NEXT

- -compare value @ <IX> with \$F0
- -if bits 7-4 are '0', then branch to NEXT

else, advance to next instr.

STATE

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Instruction Timing

• All ASM instructions require a specific amount of time to execute

-can use this information for program timing, e.g., delay loops

• $T_{instr} \equiv$ Instruction execution time

$\mathbf{T}_{\text{instr}} = \mathbf{N} \bullet \mathbf{t}_{\text{cycle}}$

- $-N \equiv #$ clock cycles per instruction
 - listed under ~* in textbook Appendix A (HO_4)
 - N depends on address mode
 - EXAMPLE:

 $-t_{cycle} \equiv$ time of one clock cycle = $1/f_{clk}$, where $f_{clk} \equiv$ clock frequency

• EXAMPLE:



Creating Time Delays

• How can a program generate a specific time delay?

- 1. Use a hardware timer -will be discussed later
- 2. Construct a <u>software delay loop</u>
- Software delay loop

-EXAMPLE:

-What is the loop delay?



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Creating Time Delays II

-EXAMPLE continued:

-What is the loop delay?

-What is 'count' for $T_D = 2ms$?

-What is the maximum delay for this loop code @ $f_{clk} = 10MHz$?



